
ON THE EXISTENCE OF MONGE MAPS FOR THE GROMOV–WASSERSTEIN PROBLEM

Joint work with T. Dumont and F.-X. Vialard.

Journée Transport Optimal et ses
applications (GdR IASIS)
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THE GROMOV–WASSERSTEIN PROBLEM

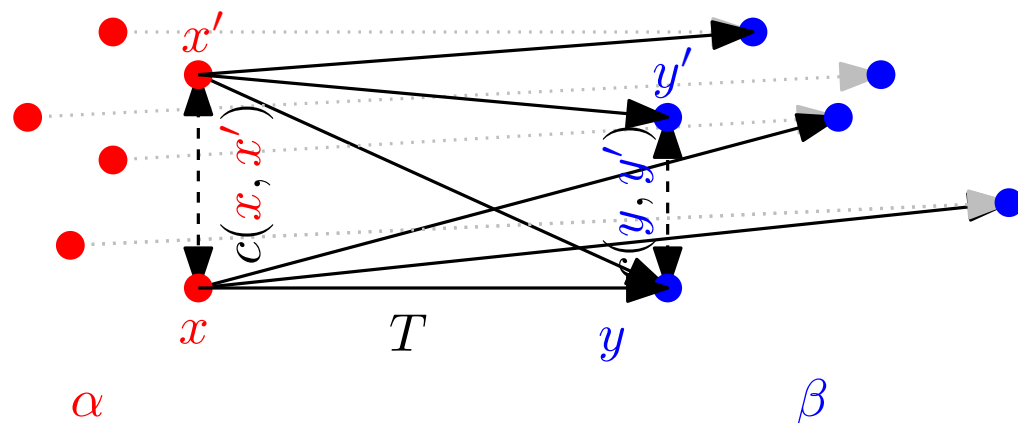
Definition:

Let $X, Y \subset \mathbb{R}^d$, $c : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}_+$ be a **cost function**, and α, β be two probability measures supported on X, Y respectively. The **Gromov–Wasserstein** problem is defined as

$$\text{GW}(\alpha, \beta) = \inf_{\pi \in \Pi(\alpha, \beta)} \iint |c(x, x') - c(y, y')|^2 d\pi(x, y) d\pi(x', y'), \quad (\text{GW})$$

where $\Pi(\alpha, \beta)$ denote the set of **transport plans** between α and β , i.e. measures on $\mathbb{R}^d \times \mathbb{R}^d$ with marginals α and β .

Example: $\alpha = \frac{1}{n} \sum_{i=1}^n \delta_{x_i}$ and $\beta = \frac{1}{n} \sum_{j=1}^n \delta_{y_j}$ ("point clouds"). Then $\pi \simeq$ a bistochastic matrix of size $n \times n$ ("bipartite graph matching"), i.e. $\pi_{ij} \leftrightarrow$ the mass transported from x_i to y_j



In a nutshell:

GW measures an **"isometry" defect** (for the cost c) between X, α and Y, β .

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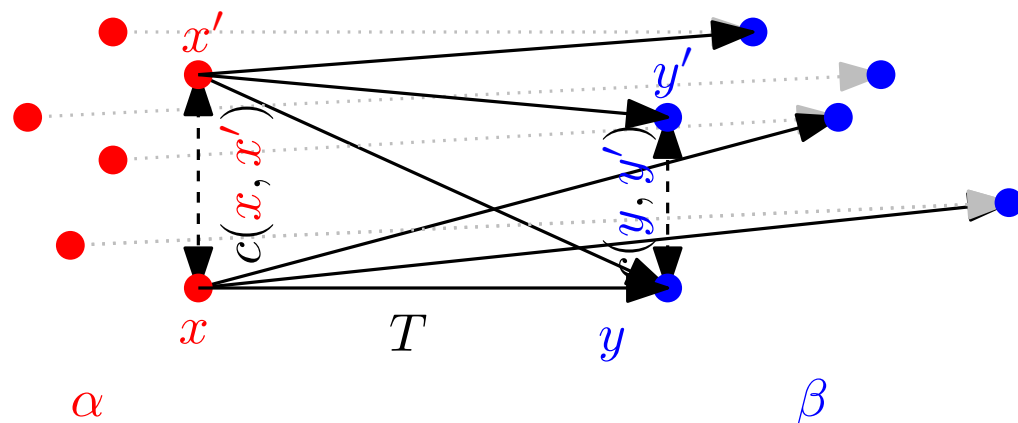
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Question 1: Under which conditions are the solution of (GW) *deterministic*, that is of the form $(\text{id}, T)_{\#} \alpha$ for some measurable map T ?

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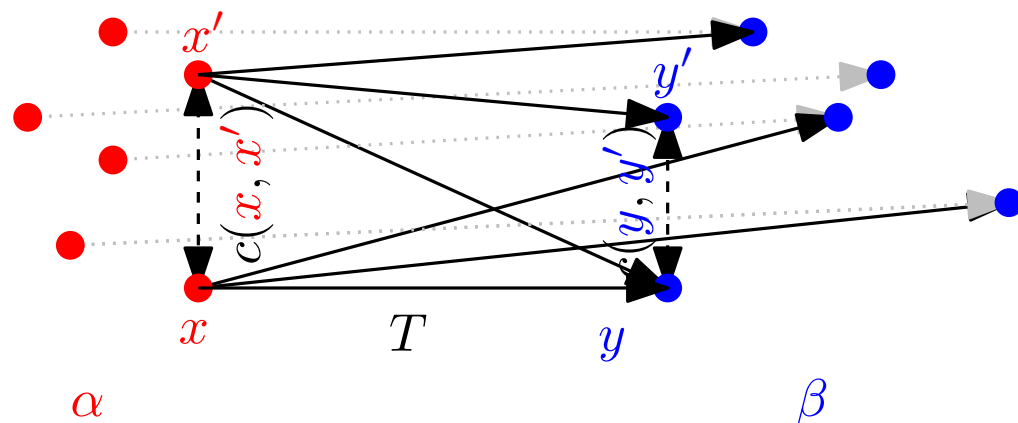
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Question 1: Under which conditions are the solution of (GW) *deterministic*, that is of the form $(\text{id}, T)_{\#} \alpha$ for some measurable map T ?

Question 2: Are there instances where GW is easy to compute? e.g. if $d = 1$?

OPTIMAL TRANSPORT MAPS FOR GW

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Proposition: [Konno 1976, Séjourné et al. 2021]

Let $c = |\cdot - \cdot|^2$ or $c = \langle \cdot, \cdot \rangle$. Let π^* solve $\text{GW}(\alpha, \beta)$. Then, any $\tilde{\pi}^*$ minimizes $\text{GW}(\alpha, \beta)$ **if and only if** it minimizes

$$\Pi(\alpha, \beta) \ni \pi \mapsto \int_{x, y} \underbrace{\left(\int_{x', y'} |c(x, x') - c(y, y')|^2 d\pi^*(x', y') \right)}_{=: C_{\pi^*}(x, y)} d\pi(x, y) = \langle C_{\pi^*}, \pi \rangle. \quad (\text{OT})$$

Idea of proof: In general, introduce the bilinear relaxation $F : (\pi, \gamma) \mapsto \iint k d\pi \otimes \gamma = \langle \pi, k\gamma \rangle$, with k symmetric, bilinear and negative on $\text{span}(\{\pi - \gamma, \pi, \gamma \in \Pi(\alpha, \beta)\})$.

Then, let (π^*, γ^*) minimize F (in particular, γ^* minimizes $\pi \mapsto F(\pi^*, \pi)$), and observe that it yields $\langle \pi^* - \gamma^*, k(\pi^* - \gamma^*) \rangle \geq 0$, hence $= 0$, hence $(\pi^* - \gamma^*) \in \ker(k)$, and thus $F(\pi^*, \pi^*) = F(\pi^*, \gamma^*)$.

It means that if π^* minimizes $\pi \mapsto F(\pi, \pi)$, it also minimizes $F(\pi^*, \pi)$, and vice-versa.

It turns out that $k(x, x', y, y') = |c(x, x') - c(y, y')|^2$ is indeed negative (on measures of mass 0) for the two costs considered in this work.

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Observation: This problem is **linear** in π . It is the **Optimal Transportation** problem between α and β for the (unknown...) cost C_{π^*} . Thanks to [Brenier 1987, Villani 2008, McCann 2011, Moameni 2016], we know that if α has a Lebesgue density, we know sufficient conditions on the cost C_{π^*} to ensure that minimizers of (OT) are induced by transport maps. Namely:

Proposition: [Moameni 2016]

If α has a density and $C : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}$ satisfies

$$\forall x_0, y_0, |\{y \mid \nabla_x C(x_0, y) = \nabla_x C(x_0, y_0)\}| \leq m, \quad (m\text{-twist})$$

then any minimizer of $\pi \mapsto \langle C, \pi \rangle$ in $\Pi(\alpha, \beta)$ is supported on the union of (at most) m transport maps.

In a nutshell:

If $\nabla_x C$ is “ m -injective”, each x is sent to (at most) m y .

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Answer: **No** (in general).

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Inner-product case: if $c = \langle \cdot, \cdot \rangle$, observe that

$$(\text{GW}) \Leftrightarrow \min_{\pi} - \iint \langle x, x' \rangle \langle y, y' \rangle d\pi^*(x', y') d\pi(x, y) \Rightarrow \text{wlog } C_{\pi^*}(x, y) = -\langle M^* x, y \rangle$$

with $M^* := \int x' y'^T d\pi^*(x', y')$.

If M^* is of full rank, $\nabla_x C_{\pi^*}(x, y) = -(M^*)^T y$ that is injective \Rightarrow 1-twist condition $\Rightarrow \exists$ optimal transport maps!^a Otherwise...?

Idea: Up to a Singular Value Decomposition, M^* is a projection and, *on the image of the projection*, the 1-twist condition is satisfied so we get the existence of a transport map. Can we lift it?

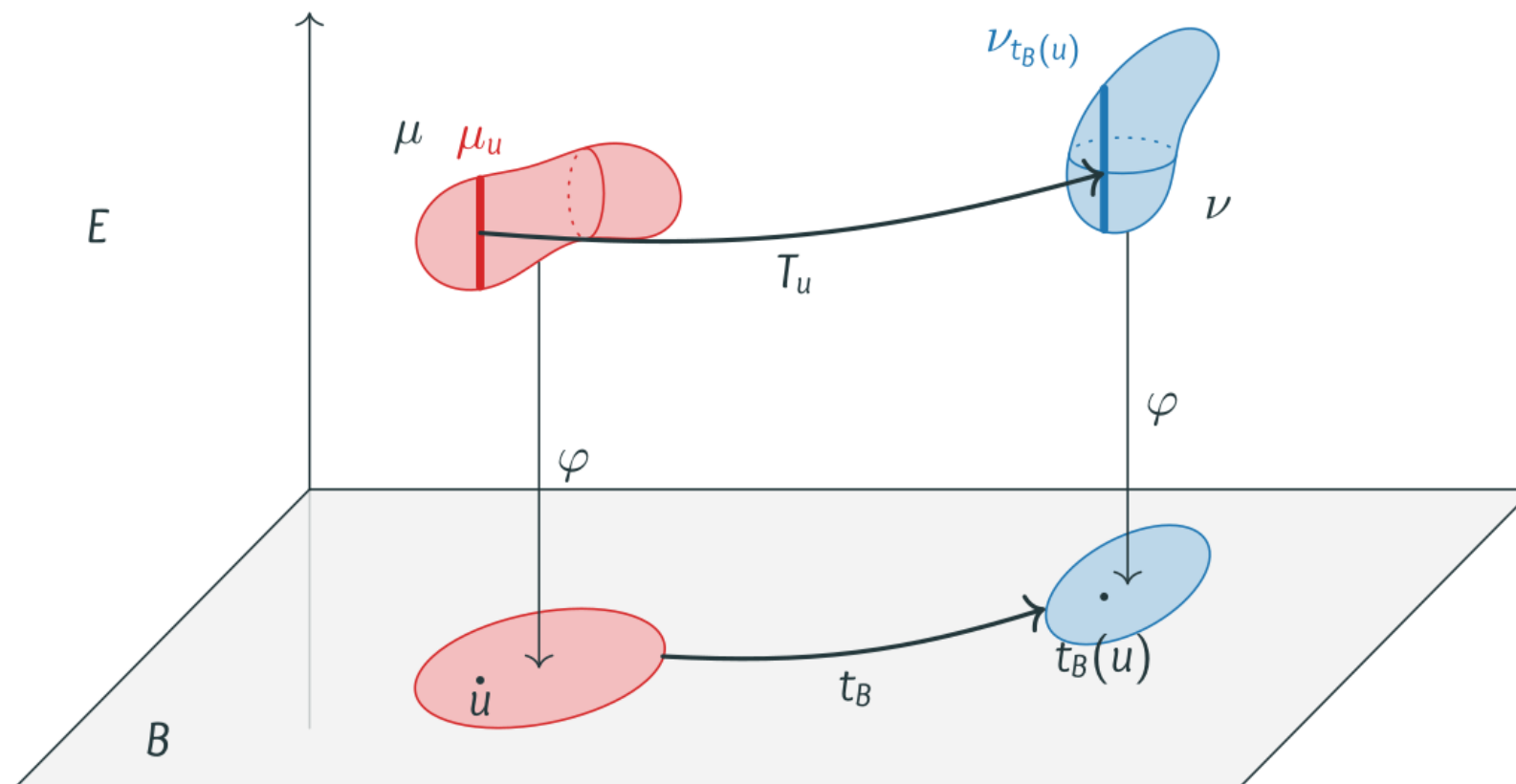
^aAlready proved by Vayer et al.

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Proposition: [Dumont, L, Vialard, 2024], “Brenier theorem for fiber-invariant costs”, Informal

Let $\alpha, \beta \in \mathcal{P}(E)$, κ a cost function, and $\varphi : E \rightarrow B$ such that (i) $\varphi\#\mu \ll \text{Vol}(B)$, (ii) $\varphi\#\alpha$ -a.e., the fiber $\varphi^{-1}(u)$ is a complete manifold (iii) $\varphi\#\alpha$ -a.e., the disintegration α_u of α wrt φ is $\ll \text{Vol}(\varphi^{-1}(u))$, (iv) $\kappa(x, y) = \tilde{\kappa}(\varphi(x), \varphi(y))$ with $\tilde{\kappa}$ twisted on B . Then \exists an optimal transport map between α and β (with some “gradient of convex function” structure on t_B and all $(T_u)_u$).



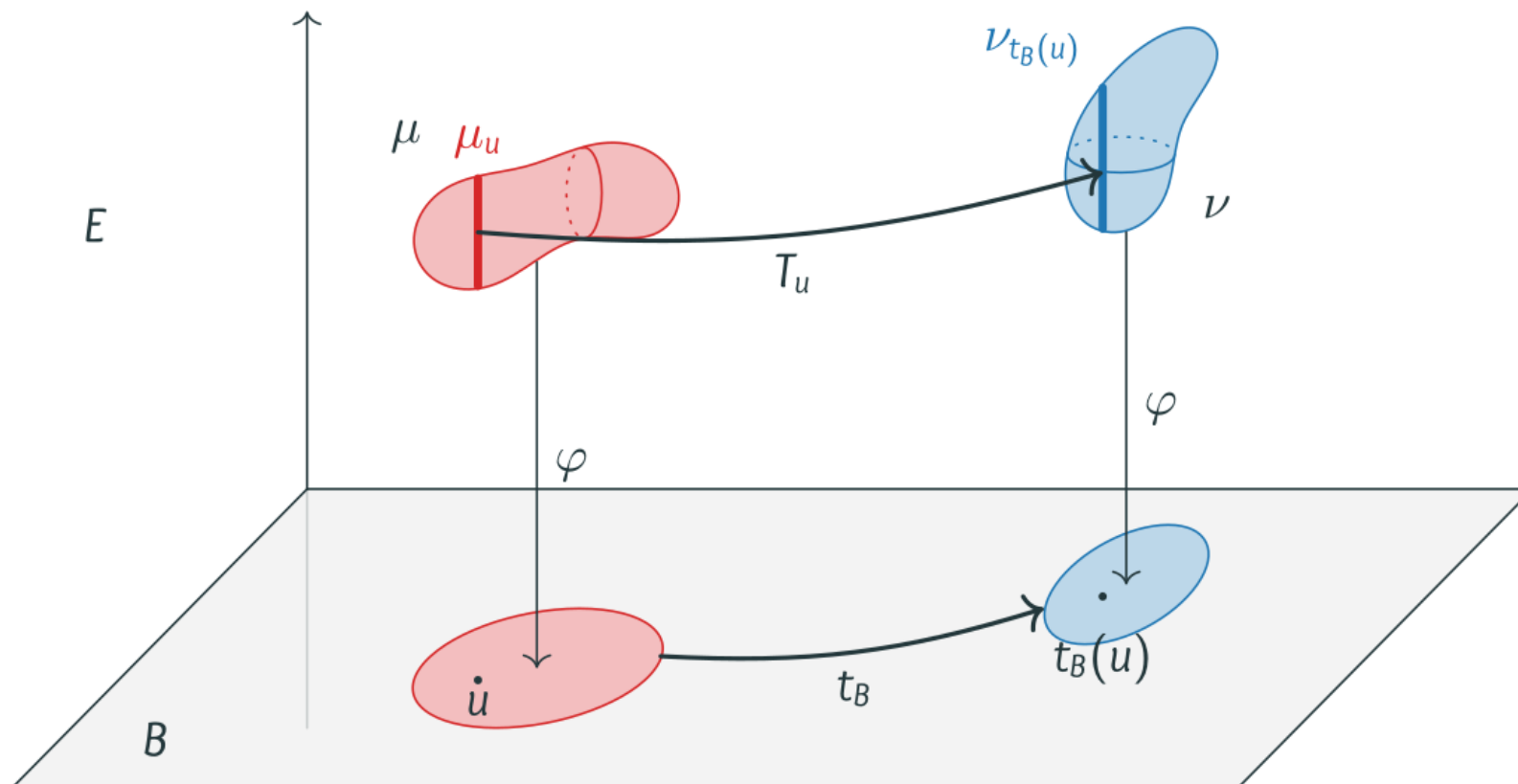
Remark: We have to ensure that the glued map $(u, x) \mapsto T_u(x)$ is measurable.
 → Adapt a result of “measurable selection of transport plans” of Fontbona et al. 2010.

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Example: $E = \mathbb{R}^d$, $\kappa(x, y) := (|x| - |y|)^2$, then take $\varphi : x \mapsto |x|$.

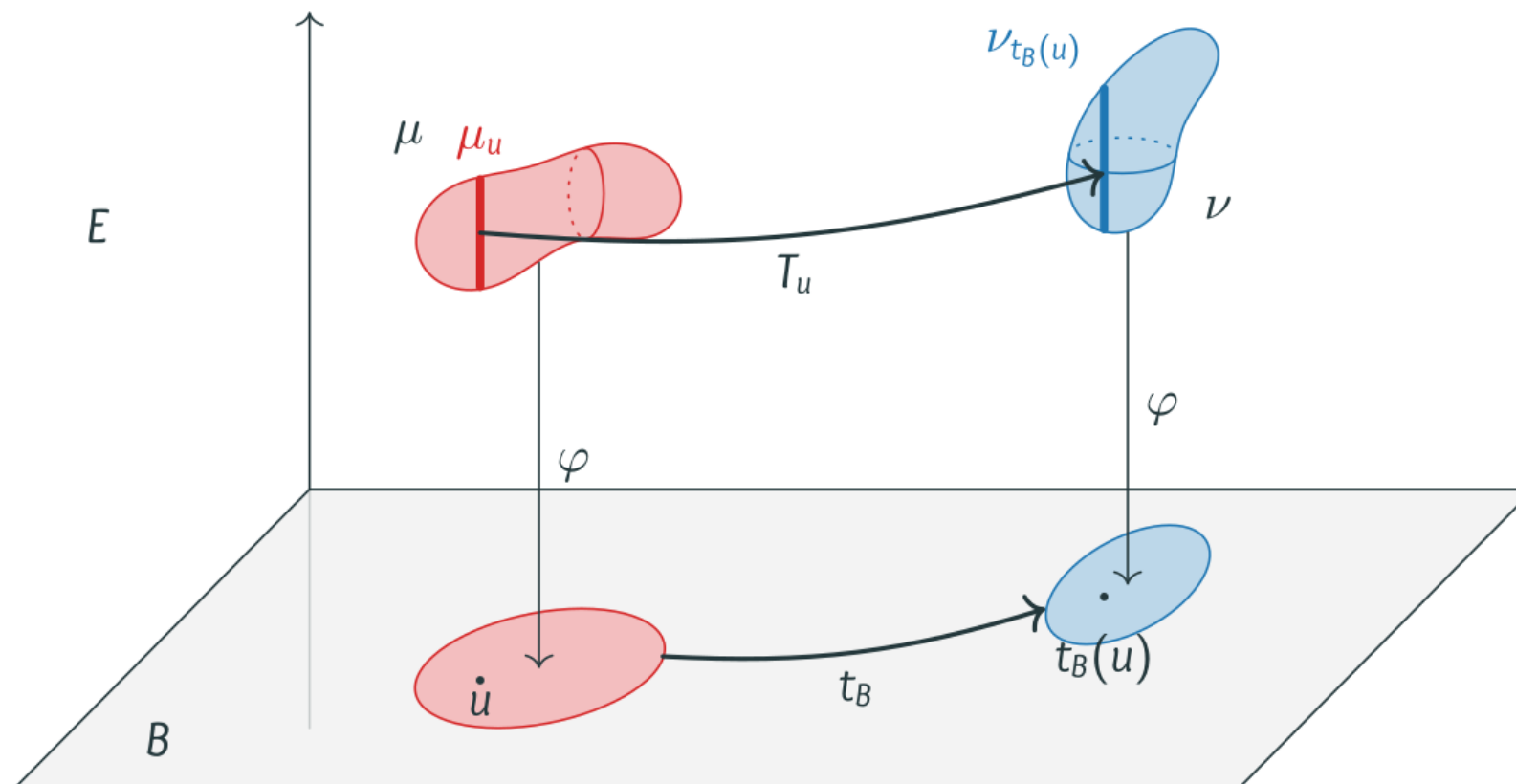
Fibers are spheres (except at $|x| = 0$ but we work a.e.), $\tilde{\kappa}$ is the quadratic cost on $B = \mathbb{R} \Rightarrow$ you can prove existence of maps for this cost.

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Corollary:

There always exist optimal transport maps for the Gromov–Wasserstein problem with inner-product cost.

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Quadratic case: Assume now that $c = |\cdot - \cdot|^2$. We get $C_{\pi^*}(x, y) = -|x|^2|y|^2 - 4 \langle M^* x, y \rangle$ (still with $M^* = \int x'^T y' d\pi^*(x', y')$).

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Proposition:

If M^* is of rank d or $d - 1$, the 2-twist condition is satisfied (idea: solutions of the equation $\nabla_x c(x, y) = \nabla_x c(x, y')$ are given by the intersection of a 1D-line and a $d - 1$ -sphere), \Rightarrow there exist “2-maps”.

If M^* is of rank $\leq d - 2$...

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If M^* is of rank $\leq d - 2$... there exists an optimal transport map!

Sketch of proof: This time, the “projection” φ is of the form $\varphi(x) = (x_H, |x_\perp|^2)$, with $x_H \in \mathbb{R}^{\text{rk}(M^*)}$ and $x_\perp \in \mathbb{R}^{d - \text{rk}(M^*)}$, the resulting cost $\tilde{\kappa}$ is twisted and the fibers are spheres of dimension $d - 1 - \text{rk}(M^*)$.

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Remark: This argument fails when $\text{rk}(M^*) = d - 1$ because fibers are $\{\pm x_\perp\} \rightarrow$ cannot build maps between fibers.

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Tightness? For the quadratic cost $c = |\cdot - \cdot|^2$, we proved that if $\text{rk}(M^*) \leq d - 2$, there exists an optimal Monge map, but we only have 2-maps if M^* is of higher rank... Did we miss something?

OPTIMAL TRANSPORT MAPS FOR GW

Recall: $\text{GW}(\alpha, \beta) = \inf_{\pi \in \Pi(\alpha, \beta)} \iint |c(x, x') - c(y, y')|^2 d\pi(x, y) d\pi(x', y')$.

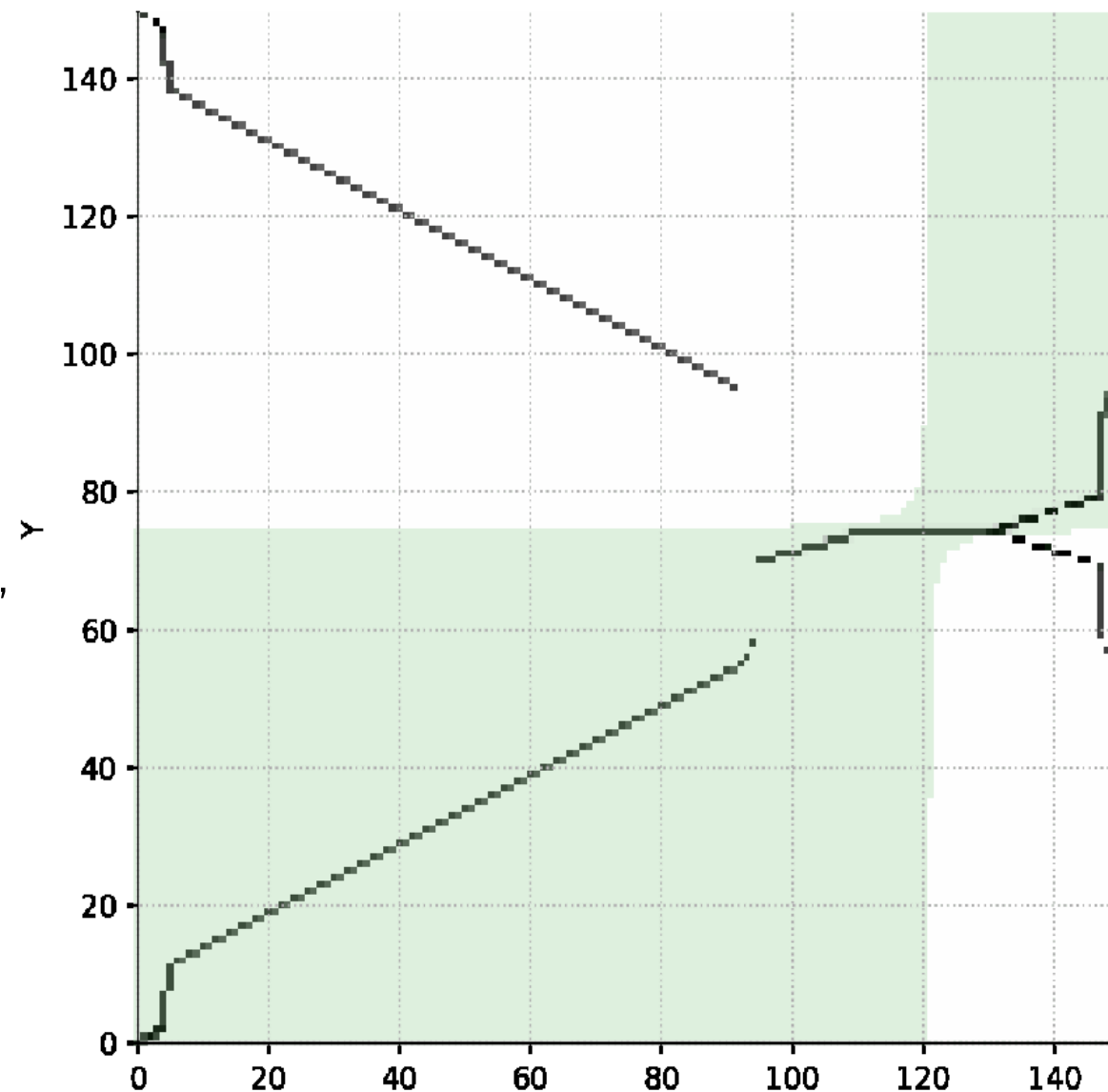
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Conjecture: No (only numerical evidence though).

More precisely, we adversarially build configurations for which the optimal GW plan looks like an actual 2-map.

Remark: Several layers of approximation/discretization :

- Discretize the ground space \mathbb{R} (starting from α with density),
- in dim 1, $M^* \in [m_{\min}, m_{\max}]$ with tractable m_{\min}/m_{\max} , and thus solve the OT problem with cost $c_m(x, y) = -x^2 y^2 - 4mxy$ for m in a discretization of this segment, yielding π_m^* , and then evaluate the GW performance of all π_m^* to find the actual optimal GW plan \rightarrow need stability results to make this rigorous.

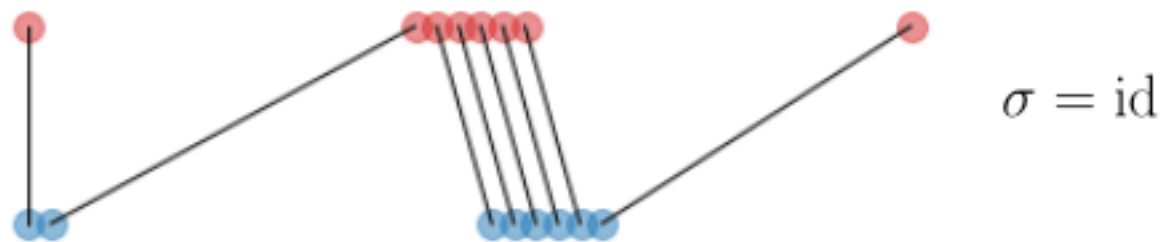


GW IN DIMENSION 1

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Observation: Assume $d = 1$, $c = |\cdot - \cdot|^2$, and let $\alpha = \frac{1}{n} \sum_{i=1}^n \delta_{x_i}$ and $\beta = \frac{1}{n} \sum_{j=1}^n \delta_{y_j}$, with $x_1 \leq \dots \leq x_n$ and $y_1 \leq \dots \leq y_n$. It would be nice if the increasing mapping $(x_i \mapsto y_i)$ or the decreasing one $(x_i \mapsto y_{n+1-i})$ would be optimal.

- Empirically satisfied, works well in applications...



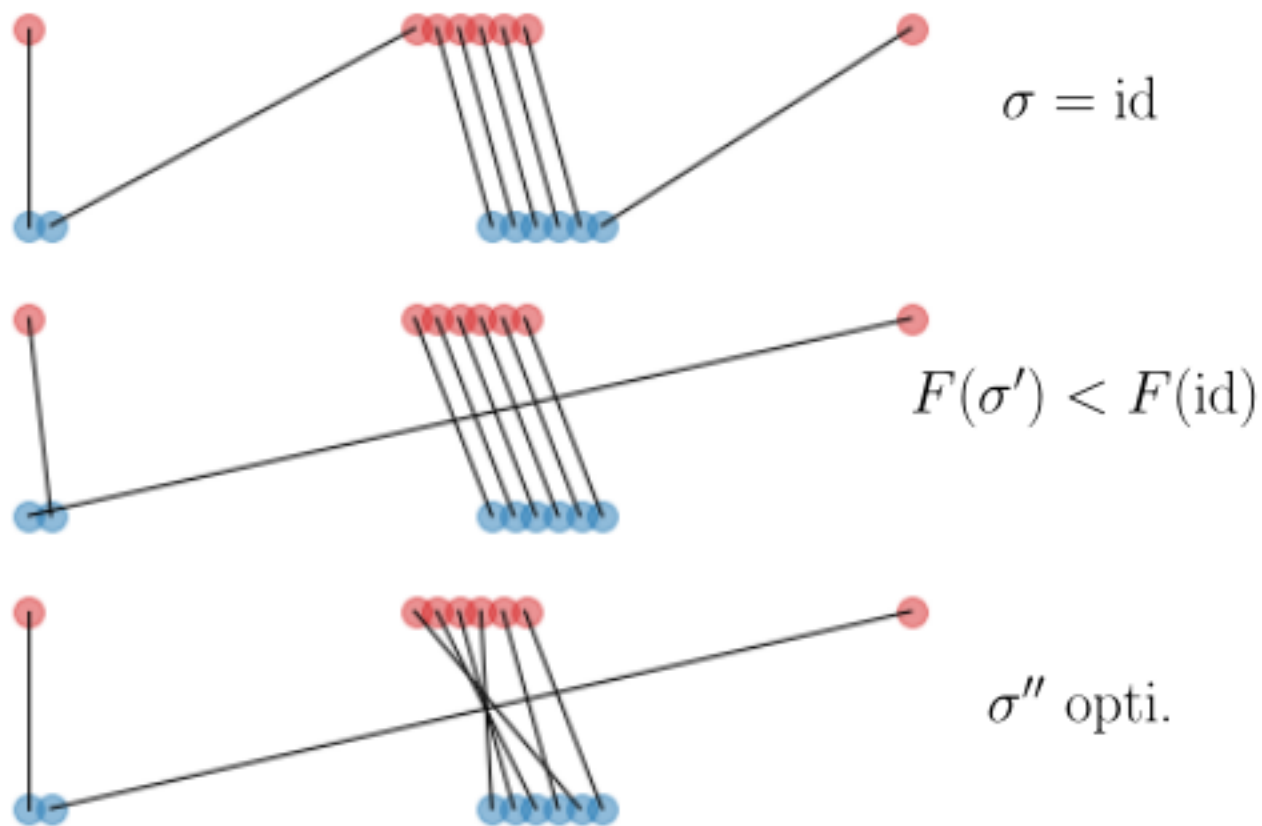
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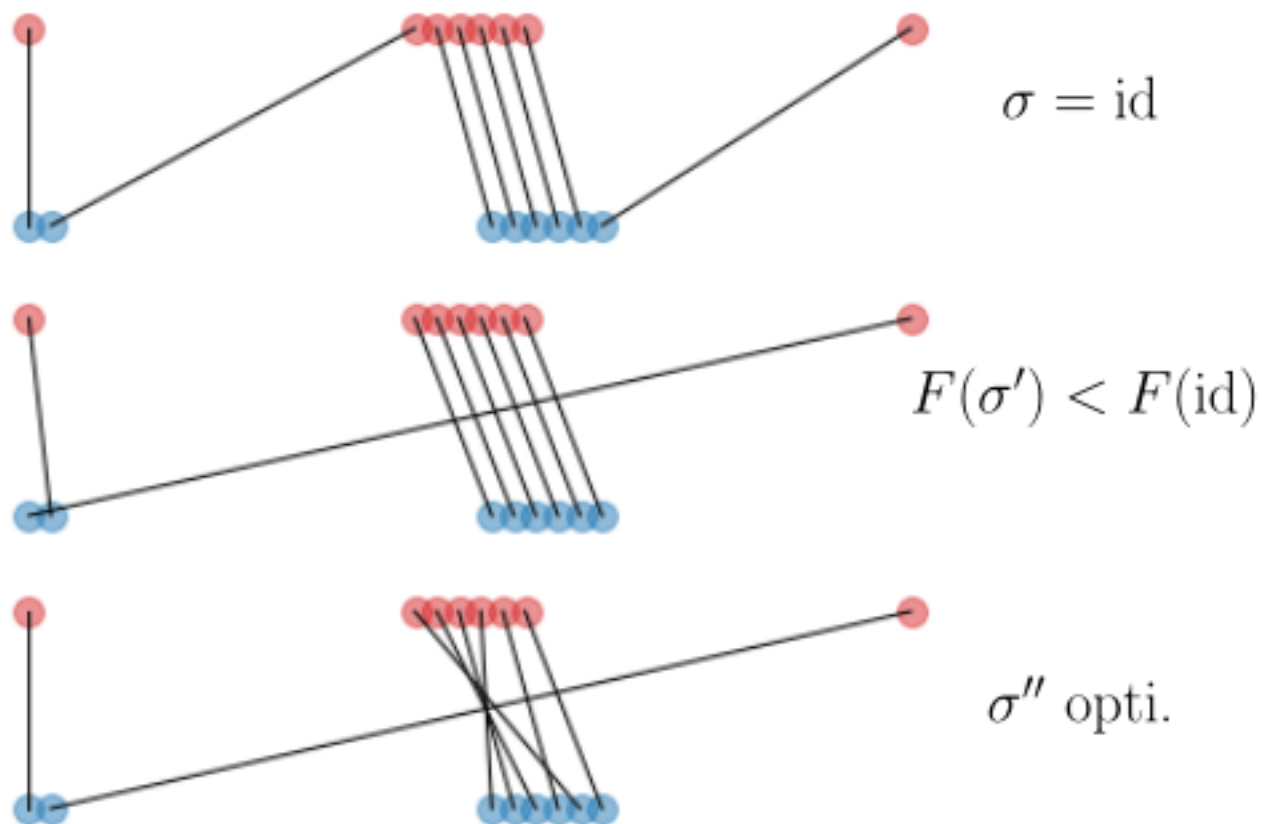
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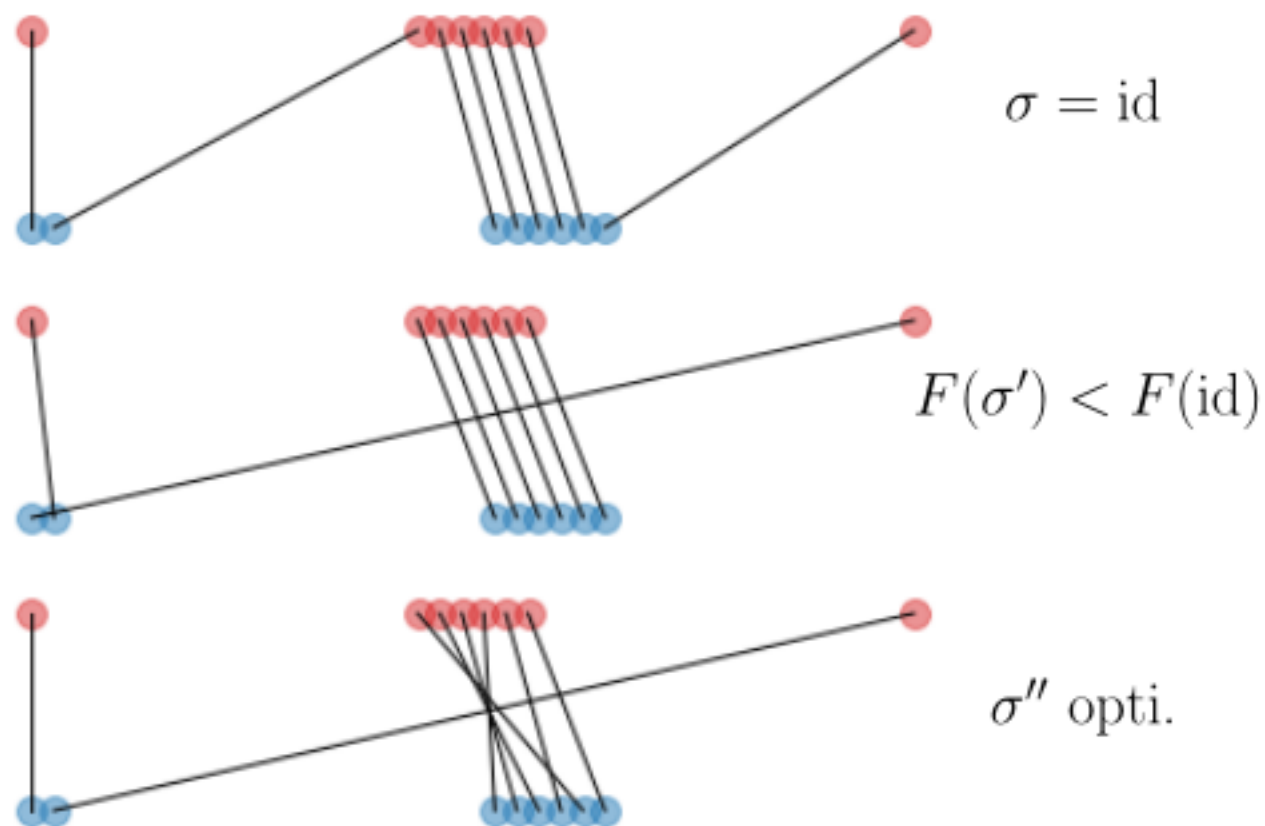
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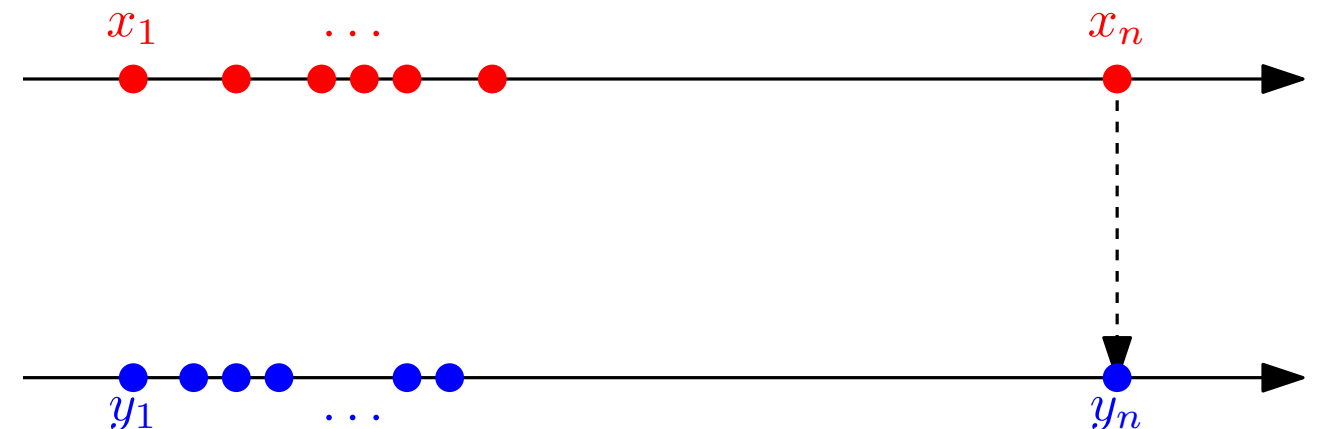
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(Open) Question: Why does it work “in practice” and only fails on “adversarial” examples?

(Partial) Answer (informal): if (say) x_n and y_n are far away from the others $(x_i, y_i)_i$, they must be matched together, and this “forces” a monotone (increasing) matching (“long-range interactions dominate”).



Conclusion & Some remaining questions

Conclusion :

- \exists (Gromov–)Monge map for GW with the inner-product cost,
- \exists 2-maps for GW with quadratic cost, and maps provided the covariance of a solution π^* is sufficiently singular. Numerical examples suggests the tightness of this result.
- (Gromov–)Monge maps are obtained by gluing a (standard) Monge map on a projection space B and Monge maps fiber-wise.
- In 1D (and probably in higher-dimension), long-range interactions dominate and dictate the structure of the matching.

Some questions left:

- Is M^* generically non-singular?
- For the quadratic cost: can we find a condition to guarantee the existence of Monge map even if M^* is of full rank? (Finding a numerical counter-example was actually quite hard!)
- Can we leverage more regularity on the disintegration $(\mu_u)_u, (\nu_u)_u$ to get more regularity/structure on the Gromov–Monge map?
- Can we leverage the existence of Gromov–Monge maps in practice?
- In 1D, can we make our observations more quantitative?

Thanks :-)